

# Community detection with the non-backtracking operator

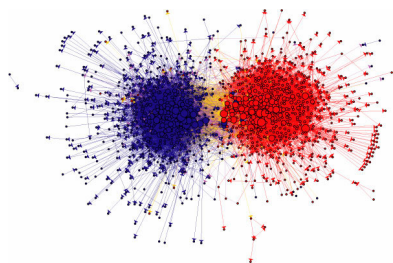
Marc Lelarge

INRIA-ENS

Aalto University, Helsinki, October 2016

# Motivation

- Community detection in social or biological networks in the sparse regime with a small average degree.

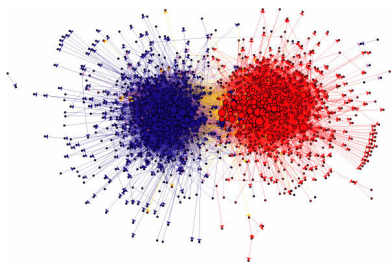


## Adamic Glance '05

- Performance analysis of spectral algorithms on a toy model (where the ground truth is known!).

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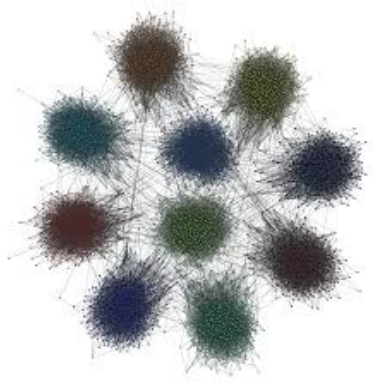
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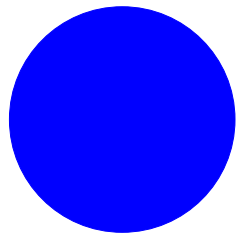
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# A model: the stochastic block model



# The sparse stochastic block model

A random graph model on  $n$  nodes with three parameters,  
 $a, b, c \geq 0$ .

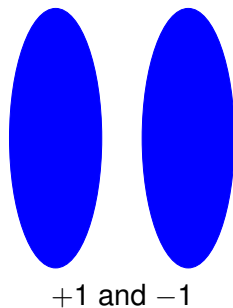


total population

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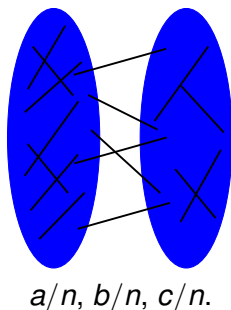
- Assign each vertex spin  $+1$  or  $-1$  uniformly at random.



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A random graph model on  $n$  nodes with three parameters,  $a, b, c \geq 0$ .

- Independently for each pair  $(u, v)$ :
  - if  $\sigma_u = \sigma_v = +1$ , draw the edge w.p.  $a/n$ .
  - if  $\sigma_u \neq \sigma_v$ , draw the edge w.p.  $b/n$ .
  - if  $\sigma_u = \sigma_v = -1$ , draw the edge w.p.  $c/n$ .



# Community detection problem

- Reconstruct the underlying communities (i.e. spin configuration  $\sigma$ ) based on one realization of the graph.
- **Asymptotics**:  $n \rightarrow \infty$
- **Sparse graph**: the parameters  $a, b, c$  are fixed.
- notion of **performance**:  
w.h.p. strictly less than half of the vertices are misclassified  
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# A first attempt: looking at degrees

- Degree in community +1 is:

$$D_+ \sim \text{Bin}\left(\frac{n}{2} - 1, \frac{a}{n}\right) + \text{Bin}\left(\frac{n}{2}, \frac{b}{n}\right)$$

- We have

$$\mathbb{E}[D_+] \approx \frac{a+b}{2}, \text{ and } \text{Var}(D_+) \approx \frac{a+b}{2}.$$

and similarly, in community -1:

$$\mathbb{E}[D_-] \approx \frac{c+b}{2}, \text{ and } \text{Var}(D_-) \approx \frac{c+b}{2}.$$

- Clustering based on degrees should 'work' as soon as:

$$(\mathbb{E}[D_+] - \mathbb{E}[D_-])^2 \succ \max(\text{Var}(D_+), \text{Var}(D_-))$$

i.e. (ignoring constant factors)

$$(a - c)^2 \succ b + \max(a, c).$$

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# Is it any good?

Data:  $A$  the adjacency matrix of the graph.

We define the mean column for each community:

$$A_+ = \frac{1}{n} \begin{pmatrix} a \\ \vdots \\ a \\ b \\ \vdots \\ b \end{pmatrix}, \text{ and } A_- = \frac{1}{n} \begin{pmatrix} b \\ \vdots \\ b \\ c \\ \vdots \\ c \end{pmatrix}$$

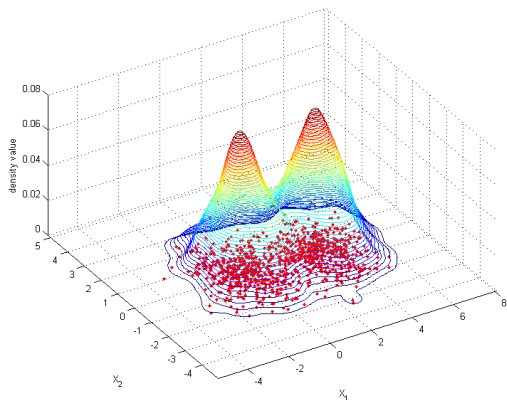
The variance of each entry is  $\leq \max(a, b, c)/n$ .

Pretend the columns are i.i.d., spherical Gaussian and  $k = n...$

# Clustering a mixture of Gaussians

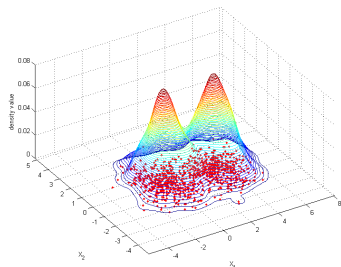
Consider a mixture of two spherical Gaussians in  $\mathbb{R}^n$  with respective means  $\mathbf{m}_1$  and  $\mathbf{m}_2$  and variance  $\sigma^2$ .

Pb: given  $k$  samples  $\sim 1/2\mathcal{N}(\mathbf{m}_1, \sigma^2) + 1/2\mathcal{N}(\mathbf{m}_2, \sigma^2)$ , recover the unknown parameters  $\mathbf{m}_1$ ,  $\mathbf{m}_2$  and  $\sigma^2$ .





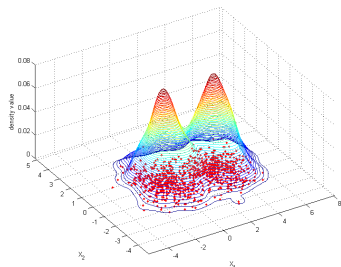
# Doing better than naive algorithm



If  $\|\mathbf{m}_1 - \mathbf{m}_2\|^2 \succ n\sigma^2$ , then the densities 'do not overlap' in  $\mathbb{R}^n$ .

Projection preserves variance  $\sigma^2$ . So projecting onto the line formed by  $\mathbf{m}_1$  and  $\mathbf{m}_2$  gives 1-dim. Gaussian variables with no overlap as soon as  $\|\mathbf{m}_1 - \mathbf{m}_2\|^2 \succ \sigma^2$ . We gain a factor of  $n$ .

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# Algorithm for clustering a mixture of Gaussians

Each sample is a column of the following matrix:

$$A = (\mathbf{A}_1, \mathbf{A}_2, \dots, \mathbf{A}_k) \in \mathbb{R}^{n \times k}$$

Consider the SVD of  $A$ :

$$A = \sum_{i=1}^n \lambda_i \mathbf{u}_i \mathbf{v}_i^T, \quad \mathbf{u}_i \in \mathbb{R}^n, \mathbf{v}_i \in \mathbb{R}^k, \lambda_1 \geq \lambda_2 \geq \dots$$

Then the best approximation for the direction  $(\mathbf{m}_1, \mathbf{m}_2)$  given by the data is  $\mathbf{u}_1$ .

Project the points from  $\mathbb{R}^n$  onto this line and then do clustering. Provided  $k$  is large enough, this 'works' as soon as:

$$\|\mathbf{m}_1 - \mathbf{m}_2\|^2 \succ \sigma^2.$$

# Back to our clustering problem

Data:  $A$  the adjacency matrix of the graph.  
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The variance of each entry is  $\leq \max(a, b, c)/n$ .

# Heuristics for community detection

The naive algorithm should work as soon as

$$\|A_+ - A_-\|^2 \succ n \underbrace{\frac{\max(a, b, c)}{n}}_{\text{Var}}$$
$$(a - b)^2 + (b - c)^2 \succ n \max(a, b, c)$$

Spectral clustering should allow you a gain of  $n$ , i.e.

$$(a - b)^2 + (b - c)^2 \succ \max(a, b, c)$$

Our previous analysis shows that clustering based on degrees works as soon as

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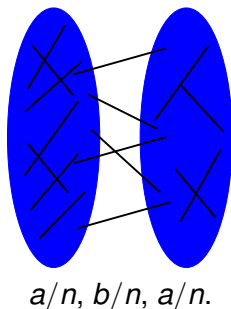
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# The sparse symmetric stochastic block model

A random graph model on  $n$  nodes with two parameters,  $a, b \geq 0$ .

- Independently for each pair  $(u, v)$ :
  - if  $\sigma_u = \sigma_v$ , draw the edge w.p.  $a/n$ .
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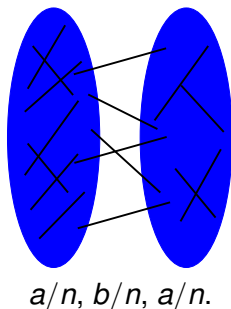
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# Efficiency of Spectral Algorithms

Boppana '87, Condon, Karp '01, Carson, Impagliazzo '01, McSherry '01, Kannan, Vempala, Vetta '04...

## Theorem

*Suppose that for sufficiently large  $K$  and  $K'$ ,*

$$\frac{(a-b)^2}{a+b} \geq (\succ)K + K' \ln(a+b),$$

*then 'trimming+spectral+greedy improvement' outputs a positively correlated (almost exact) partition w.h.p.*

Coja-Oghlan '10

Heuristic based on analogy with mixture of Gaussians:

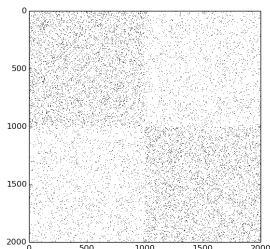
$$(a-b)^2 \succ a+b$$

# Another look at spectral algorithms

Take a finite, simple, non-oriented graph  $G = (V, E)$ .

Adjacency matrix : symmetric, indexed on vertices, for  $u, v \in V$ ,

$$A_{uv} = 1(\{u, v\} \in E).$$



Low rank approximation of the adjacency matrix works as soon as

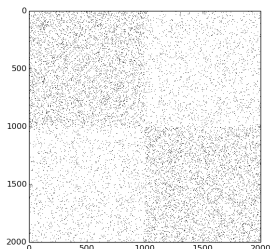
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# Spectral analysis

Assume that  $a \rightarrow \infty$ , and  $a - b \approx \sqrt{a + b}$  so that  $a \sim b$ .

$$A = \frac{a+b}{2} \frac{\mathbf{1} \mathbf{1}^T}{\sqrt{n} \sqrt{n}} + \frac{a-b}{2} \frac{\sigma \sigma^T}{\sqrt{n} \sqrt{n}} + A - \mathbb{E}[A]$$

$\frac{a+b}{2}$  is the **mean degree** and degrees in the graph are very concentrated if  $a \succ \ln n$ . We can construct

$$A - \frac{a+b}{2n} J = \frac{a-b}{2} \frac{\sigma \sigma^T}{\sqrt{n} \sqrt{n}} + A - \mathbb{E}[A]$$

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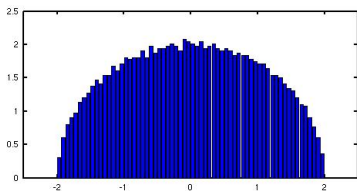
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# Spectrum of the noise matrix

The matrix  $A - \mathbb{E}[A]$  is a symmetric random matrix with independent centered entries having variance  $\sim \frac{a}{n}$ . To have convergence to the **Wigner semicircle law**, we need to normalize the variance to  $\frac{1}{n}$ .



$$ESD\left(\frac{A - \mathbb{E}[A]}{\sqrt{a}}\right) \rightarrow \mu_{sc}(x) = \begin{cases} \frac{1}{2\pi} \sqrt{4 - x^2}, & \text{if } |x| \leq 2; \\ 0, & \text{otherwise.} \end{cases}$$

# Naive spectral analysis

To sum up, we can construct:

$$\begin{aligned} M &= \frac{1}{\sqrt{a}} \left( A - \frac{a+b}{2n} J \right) \\ &= \theta \frac{\sigma}{\sqrt{n}} \frac{\sigma^T}{\sqrt{n}} + \frac{A - \mathbb{E}[A]}{\sqrt{a}}, \end{aligned}$$

with  $\theta = \frac{a-b}{\sqrt{2(a+b)}}$ .

We should be able to detect signal as soon as

$$\theta > 2 \Leftrightarrow \frac{(a-b)^2}{2(a+b)} > 4$$



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# We can do better!

A lower bound on the spectral radius of  $M = \theta \frac{\sigma}{\sqrt{n}} \frac{\sigma^T}{\sqrt{n}} + W$ :

$$\lambda_1(M) = \sup_{\|x\|=1} \|Mx\| \geq \left\| M \frac{\sigma}{\sqrt{n}} \right\|$$

But

$$\begin{aligned} \left\| M \frac{\sigma}{\sqrt{n}} \right\|^2 &= \theta^2 + \left\| W \frac{\sigma}{\sqrt{n}} \right\|^2 + 2 \langle W, \frac{\sigma}{\sqrt{n}} \rangle \\ &\approx \theta^2 + \frac{1}{n} \sum_{i,j} W_{ij}^2 \\ &\approx \theta^2 + 1. \end{aligned}$$

As a result, we get

$$\lambda_1(M) > 2 \Leftrightarrow \theta > 1 \Leftrightarrow (a-b)^2 > 2(a+b).$$

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# Baik, Ben Arous, P ech  phase transition

Rank one perturbation of a Wigner matrix:

$$\lambda_1(\theta\sigma\sigma^T + W) \xrightarrow{\text{a.s.}} \begin{cases} \theta + \frac{1}{\theta} & \text{if } \theta > 1, \\ 2 & \text{otherwise.} \end{cases}$$

Let  $\tilde{\sigma}$  be the eigenvector associated with  $\lambda_1(\theta\sigma\sigma^T + W)$ , then

$$|\langle \tilde{\sigma}, \sigma \rangle|^2 \xrightarrow{\text{a.s.}} \begin{cases} 1 - \frac{1}{\theta^2} & \text{if } \theta > 1, \\ 0 & \text{otherwise.} \end{cases}$$

Watkin Nadal '94, Baik, Ben Arous, P ech  '05  
Newman, Rao '14

For SBM with  $a, b \rightarrow \infty$ ,

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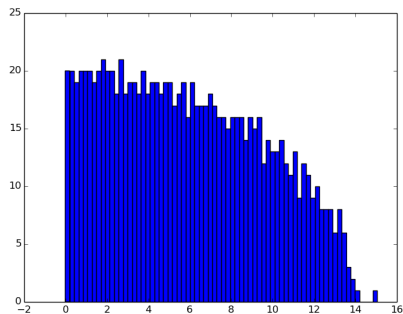
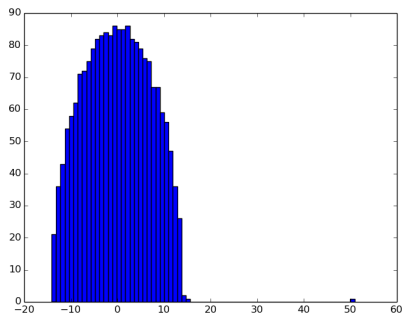
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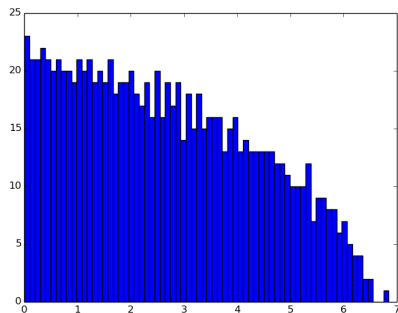
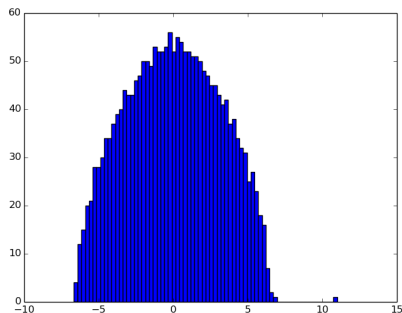
# When $a, b \rightarrow \infty$ spectral is optimal



SBM with  $n = 2000$ , average degree 50 and  $\frac{(a-b)^2}{2(a+b)} = 2$ .

Random matrix theory predicts  $\lambda_1 = 51$ ,  $\lambda_2 = 15$  and noise at  $|\lambda_3| < 14.14$

# Decreasing the average degree

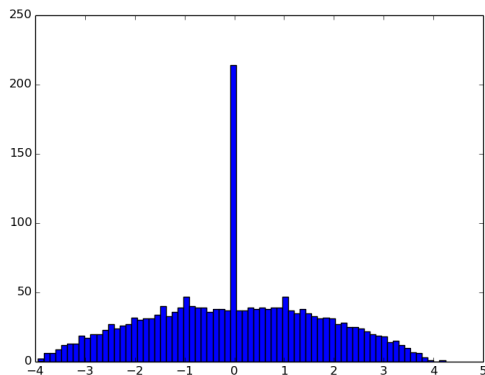


SBM with  $n = 2000$ , average degree 10 and  $\frac{(a-b)^2}{2(a+b)} = 2$ .

Random matrix theory predicts  $\lambda_1 = 11$ ,  $\lambda_2 = 6.7$  and noise at  $|\lambda_3| < 6.3$



# Problems when the average degree is small



SBM with  $n = 2000$ , average degree 3 and  $\frac{(a-b)^2}{2(a+b)} = 2$ .

Random matrix theory predicts  $\lambda_1 = 4$ ,  $\lambda_2 = 3.67$  and noise at  $|\lambda_3| < 3.46$

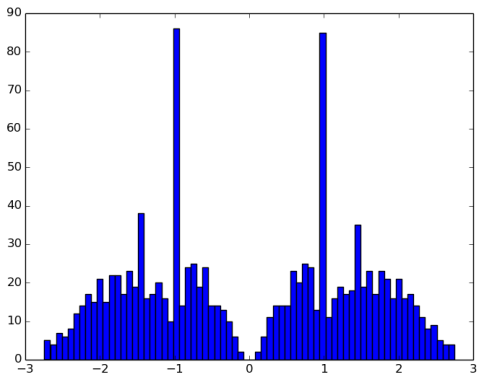
# Problems when the average degree is finite

- **High degree nodes:** a star with degree  $d$  has eigenvalues  $\{-\sqrt{d}, 0, \sqrt{d}\}$ .  
In the regime where  $a$  and  $b$  are finite, the degrees are asymptotically Poisson with mean  $\frac{a+b}{2}$ . The adjacency matrix has  $\Omega\left(\sqrt{\frac{\ln n}{\ln \ln n}}\right)$  eigenvalues.
- **Low degree nodes:** instead of the adjacency matrix, take the (normalized) Laplacian but then isolated edges produce spurious eigenvalues.

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- **High degree nodes:** a star with degree  $d$  has eigenvalues  $\{-\sqrt{d}, 0, \sqrt{d}\}$ .  
In the regime where  $a$  and  $b$  are finite, the degrees are asymptotically Poisson with mean  $\frac{a+b}{2}$ . The adjacency matrix has  $\Omega\left(\sqrt{\frac{\ln n}{\ln \ln n}}\right)$  eigenvalues.
- **Low degree nodes:** instead of the adjacency matrix, take the (normalized) Laplacian but then isolated edges produce spurious eigenvalues.

# Problems when the average degree is small



Same graph after trimming.

# Phase transition for $a, b = O(1)$

## Theorem

$$\tau = \frac{(a - b)^2}{2(a + b)}$$

*If  $\tau > 1$ , then positively correlated reconstruction is possible.*

*If  $\tau < 1$ , then positively correlated reconstruction is impossible.*

Conjectured by **Decelle, Krzakala, Moore, Zdeborova '11** based on statistical physics arguments.

- Non-reconstruction proved by **Mossel, Neeman, Sly '12**.
- Reconstruction proved by **Massoulié '13** and **Mossel, Neeman, Sly '13**.

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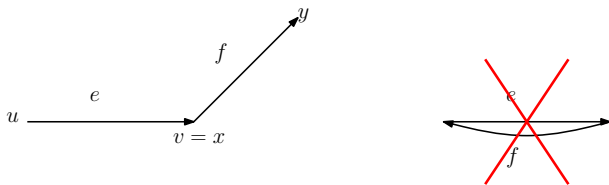
# Regularization through the non-backtracking matrix

Let  $\vec{E} = \{u \rightarrow v; \{u, v\} \in E\}$  be the set of oriented edges.

$m = |\vec{E}|$  is twice the number of unoriented edges.

The **non-backtracking matrix** is an  $m \times m$  matrix defined by

$$B_{u \rightarrow v, v \rightarrow w} = 1(\{u, v\} \in E)1(\{v, w\} \in E)1(u \neq w)$$



$B$  is NOT symmetric:  $B^T \neq B$ . We denote its eigenvalues by  $\lambda_1, \lambda_2, \dots$  with  $\lambda_1 \geq \dots \geq |\lambda_m|$ .

Proposed by Krzakala et al. '13.



# Ihara-Bass' Identity

Let  $D$  the diagonal matrix with  $D_{vv} = \deg(v)$ . We have

$$\det(z - B) = (z^2 - 1)^{|E| - |V|} \det(z^2 - Az + D - Id)$$

If  $G$  is  $d$ -regular, then  $D = dId$  and,

$$\sigma(B) = \{\pm 1\} \cup \left\{ \lambda : \lambda^2 - \lambda\mu + (d - 1) = 0 \text{ with } \mu \in \sigma(A) \right\}.$$

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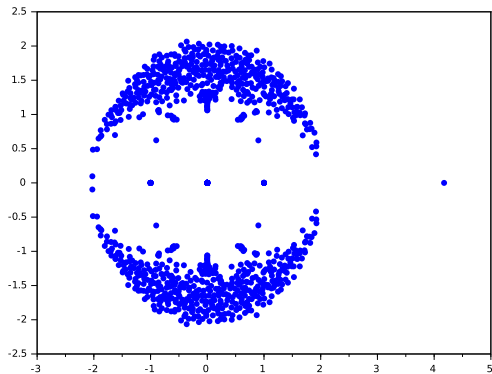
# Non-Backtracking matrix of regular graphs

For a  $d$ -regular graph,  $\lambda_1 = d - 1$ ,

- ★ Alon-Boppana bound :  $\max_{k \neq 1} \Re(\lambda_k) \geq \sqrt{\lambda_1} - o(1)$ .
- ★ Ramanujan (non bipartite) :  $|\lambda_2| = \sqrt{\lambda_1}$
- ★ Friedman's thm :  $|\lambda_2| \leq \sqrt{\lambda_1} + o(1)$  if  $G$  random uniform.

# Simulation for Erdős-Rényi Graph

Eigenvalues of  $B$  for an Erdős-Rényi graph  $G(n, \lambda/n)$  with  $n = 500$  and  $\lambda = 4$ .



# Erdős-Rényi Graph

Eigenvalues of  $B$ :  $\lambda_1 \geq |\lambda_2| \geq \dots$

## Theorem

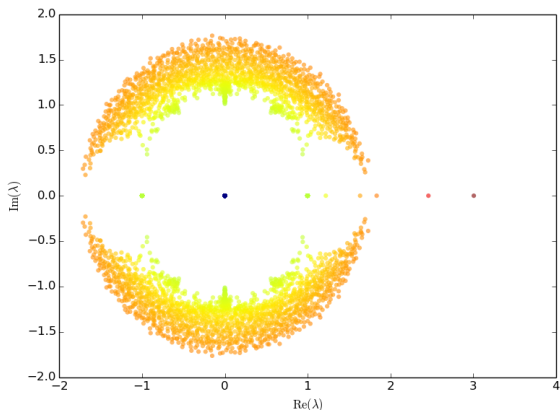
*Let  $\lambda > 1$  and  $G$  with distribution  $G(n, \lambda/n)$ . With high probability,*

$$\begin{aligned}\lambda_1 &= \lambda + o(1) \\ |\lambda_2| &\leq \sqrt{\lambda} + o(1).\end{aligned}$$

Bordenave, Lelarge, Massoulié '15

# Simulation for Stochastic Block Model

Eigenvalues of  $B$  for a Stochastic Block Model with  $n = 2000$ ,  
mean degree  $\frac{a+b}{2} = 3$  and  $\frac{a-b}{2} = 2.45$



# Stochastic Block Model

Eigenvalues of  $B$ :  $\lambda_1 \geq |\lambda_2| \geq \dots$

## Theorem

*Let  $G$  be a Stochastic Block Model with parameters  $a, b$ . If  $(a - b)^2 > 2(a + b)$ , then with high probability,*

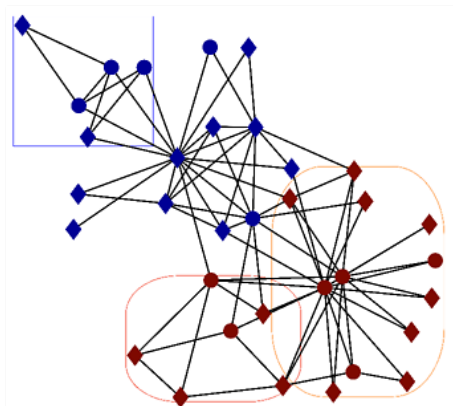
$$\lambda_1 = \frac{a + b}{2} + o(1)$$

$$\lambda_2 = \frac{a - b}{2} + o(1)$$

$$|\lambda_3| \leq \sqrt{\frac{a + b}{2}} + o(1).$$

Bordenave, Lelarge, Massoulié '15

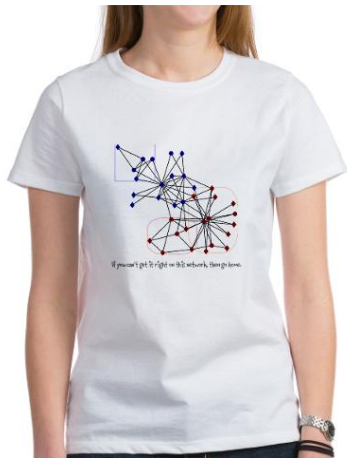
# Test with real benchmarks



If you can't get it right on this network, then go home.

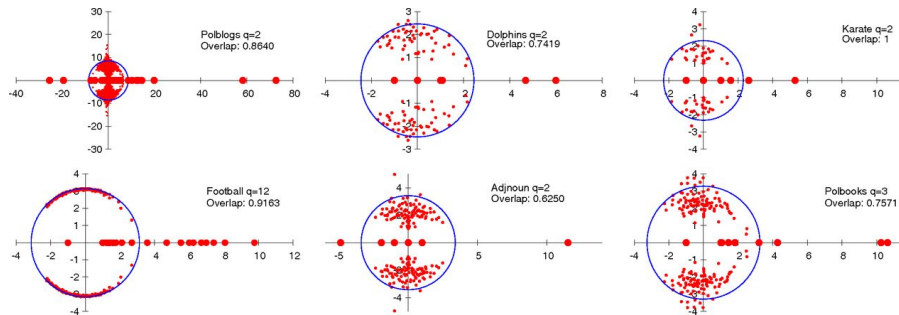


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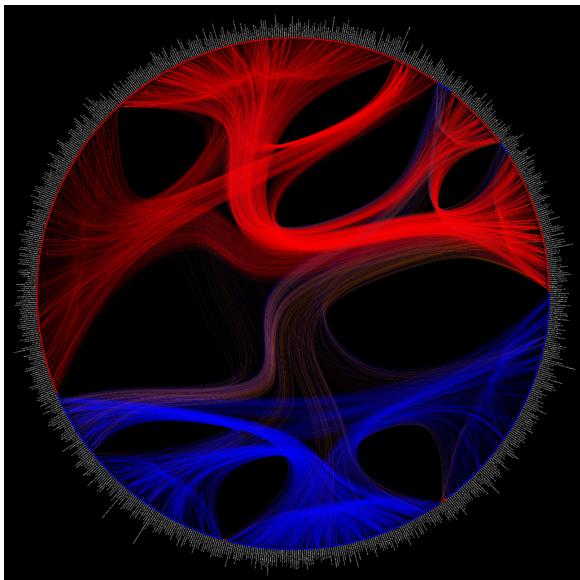
The Power Law Shop

# The non-backtracking matrix on real data



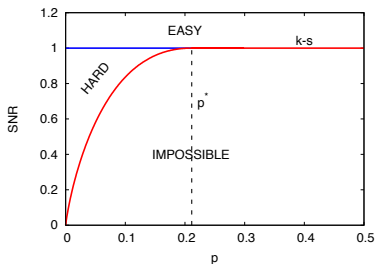
from Krzakala, Moore, Mossel, Neeman, Sly, Zdeborová '13

# Back to political blogging network data



# Non-symmetric Stochastic Block Model

Consider the case where there is a small community of size  $pn$  with  $p < 1/2$ , then the SNR is given by  $d(1 - b)^2$  where  $d$  is the average degree.

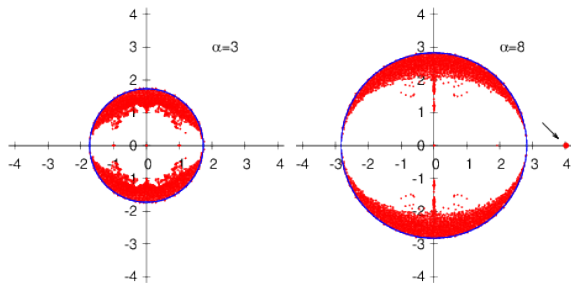


Phase diagram with  $p^* = \frac{1}{2} - \frac{1}{2\sqrt{3}}$ .

Lelarge, Caltagirone & Miolane, '16

# Some extensions

For the **labeled** stochastic block model, we also conjecture a **phase transition**. We have partial results and an optimal spectral algorithm.



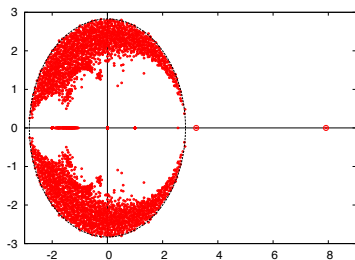
Saade, Krzakala, Lelarge, Zdeborová, '15,'16

## Some extensions

The non-backtracking matrix is also working for the **degree-corrected SBM**.

ongoing work with **Gulikers and Massoulié**.

We can adapt the non-backtracking matrix to deal with small **cliques**.

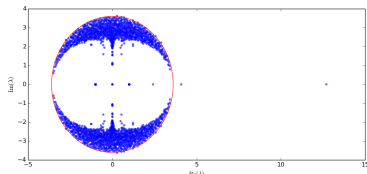
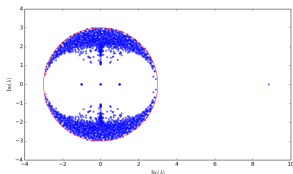


ongoing work with **Caltagirone**.

# Some extensions

SBM with no noise  $b = 0$  but with **overlap**.

Spectrum of the non-backtracking operator with  $n = 1200$ ,  
 $sn = 400$  and  $a = 9$  and  $13$ . The circle has radius  $\sqrt{a(2 - 3s)}$   
in each case.

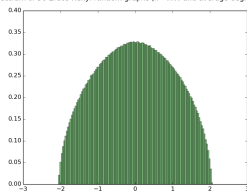


Kaufmann, Bonald, Lelarge '16

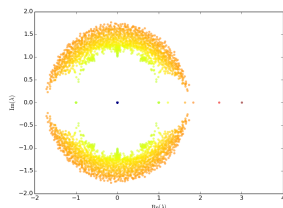
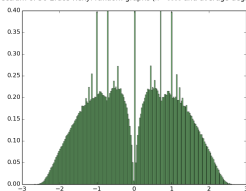
# Non-backtracking vs adjacency

On the **sparse stochastic block model** with probability of intra-edge  $a/n$  and inter-edge  $b/n$ .

Spectrum of 50 Erdos Renyi random graphs ( $n=5000$  and average degree  $c=20$ )



Spectrum of 50 Erdos Renyi random graphs ( $n=5000$  and average degree  $c=2$ )



**The problem:** if  $a, b \rightarrow \infty$ , then Wigner's semi-circle law + BBP phase transition but if  $a, b < \infty$  as  $n \rightarrow \infty$ , then **Lifshitz tails**.

**The solution:** the non-backtracking matrix on directed edges of the graph:  $B_{u \rightarrow v, v \rightarrow w} = 1(\{u, v\} \in E)1(\{v, w\} \in E)1(u \neq w)$  achieves **optimal detection** on the SBM.

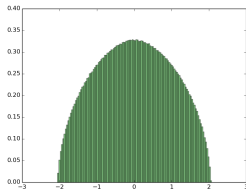
THANK YOU!



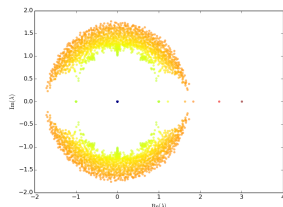
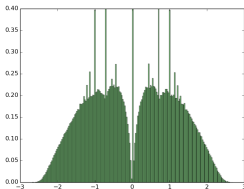
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THANK YOU!